

Facility layout for cellular manufacturing system under dynamic conditions

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CHRONICLE

ABSTRACT

Article history:

Received October 25, 2015

Received in revised format:

December 12, 2015

Accepted February 4, 2016

Available online

February 6 2016

Keywords:

Cellular manufacturing

Dynamic conditions

Facility layout

Cell formation

Hierarchical Genetic Algorithm

Cellular manufacturing is considered as a lean technique of producing similar parts using sells or groups of team members, workstations, or equipment to facilitate operations by removing setup and unnecessary cost components among various operations. Cell formation and layout planning are the most components of the cellular manufacturing. This paper presents a dynamic method to minimize different costs including the total cost of movements within and between cells and exceptional parts. In this study, the Hierarchical Genetic Algorithm (HGA) is used for solving the resulted model and the results are compared with genetic algorithm. The results have indicated that the proposed method could reach optimal solutions for some small and medium sized problems in reasonable amount of time.

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1. Introduction

Cellular manufacturing is considered as a lean technique of producing similar parts using sells or groups of team members, workstations, or equipment to facilitate operations by removing setup and unnecessary cost components among various operations (Ming & Ponnambalam, 2008). Cell formation and layout planning are the most components of the cellular manufacturing (Dixit & Mishra, 2009). Co and Araar (1988) presented a three-stage procedure for configuring machines into manufacturing cells, and assigning the cells to process specific sets of jobs and offered a direct-search method for defining the composition of manufacturing cells. Vakharia and Wemmerlov (1990) offered a cell formation technique, which combines the problems of cell formation and within-cell material flows. Houshyar and McGINNIS (1990) implemented a cut method to the corresponding quadratic assignment

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problem, and proposed a heuristic which produces a good solution. Logendran (1991) investigated the effect of sequence of operations and layout of cells in cellular manufacturing. They offered a technique, which includes two important factors and presented for determining optimal/near-optimal machine-part clusters in cellular manufacturing. Alfa et al. (1992) presented the simultaneous solution of the machine grouping and layout problems in cellular manufacturing systems. Balasubramanian and Panneerselvam (1993) investigated a method, which facilitates the formation, of machine groups. For a given set of components and their process routeing, the method first looked for several possible cells and then proposed a similarity matrix to reflect the components' routeing similarity with that of the arrangement of the facilities in the cell. They also proposed a heuristic procedure for tackling the 'covering problem' to detect an economical number of manufacturing cells and cell arrangements such that each cell is determined to process specific component(s).

Chandrasekharan and Rajagopalan (1993) proposed a multidimensional scaling algorithm for group layout in cellular manufacturing. Tavakkoli-Moghaddam et al. (2007) presented a model to solve a facility layout problem in cellular manufacturing systems (CMSs) with stochastic demands. Chan et al. (2008) proposed a two-phase approach to tackle the cell formation problem (CFP) by taking into account both intra-cell and inter-cell part movements. Mahdavi and Mahadevan (2008) proposed an algorithm for cellular manufacturing system and layout design based on sequence data. Paydar et al. (2008) proposed a fuzzy goal programming based method to handle an extended multi-objective mathematical linear programming model to tackle two essential problems in cellular manufacturing systems simultaneously: Cell formation and layout design.

Ahi et al. (2009) presented a method to determine cell formation, intracellular machine layout and cell layout in the CMS problem based on Technique for Order of Preference by Similarity to Ideal Solution (TOPSIS) (Hwang & Yoon, 1981) method. Chang et al. (2009) presented a method for cell formation and cell layout design in cellular manufacturing system. Mahdavi et al. (2010) proposed a mathematical model for the joint problem of the cell formation problem and the machine layout where the objective was to minimize the total cost of inter-cell and intra-cell movements and the investment cost of machines. Paydar et al. (2010) applied simulated annealing for designing cellular manufacturing systems. Nouri and Hong (2013) developed a bacteria foraging optimization algorithm for cell formation in cellular manufacturing system considering cell load variations. Chandrasekar and Venkumar (2013) proposed a genetic algorithm method for integrating cell formation with machine layout and cell layout.

2. The proposed study

The proposed study of this paper considers the following objectives,

1. It groups equipment and part families.
2. The proposed method reduces depreciation cost as well as the cost of transportation from one period to another period,
3. It considers minimization of the cost of transportsations of parts between and inside cells.
4. It reducing the number of exceptional parts.
5. It also reduces the number machine-movement from one cell to another one from one period to another.

Assumptions

1. Each part needs pre-specified operations and each operation has a known processing time.
2. The capacity of each machine is known in advance.
3. Parts are moved inside and between cells in batch and the cost of transportation is known.
4. There is a known maximum number of machines in each cell.

5. Any machine can be moved from one cell to another between each period.
6. The cost of transportation between cells depends on the sequence of cells.
7. Each machine is capable of performing one type of operations.
8. The cost of transportation between cells is fixed.
9. No inventory exists between periods and there is no backorders.
10. All machines are operating with 100% efficiencies.
11. There is one line transportation operations.
12. The operations of each part on each machine is known in advance.
13. All machines are leased for operations.
14. The maintenance, depreciation and other machine related cost associated with each machine are fixed and known in advance.
15. The places assigned for all machines in every cells are known in advance.

Symbols

i, i'	Index for machines	$i, i' = 1, 2, \dots, m$
j	Index for parts	$j = 1, 2, \dots, n$
p	The position of each machine	$p = 1, 2, \dots, mp$
k, k'	Index for cells	$k, k' = 1, 2, \dots, C$
h	Index for planning time	$h = 1, 2, \dots, H$

Parameters

B_j	Size of batch for movement of part j
C_{intra}^j	The cost of intra cell movement of each part j
C_{inter}^j	The cost of movement of each part j between cells
C_{back}^j	The cost of backward movement of each part j inside the cell
M_{ih}	The position of machine i on period h
R_{ij}	The number of operations for part j on machine i
D_{jh}	Demand for part j in period h
NM	Maximum number of machines in each cell
$f_{ii'h}^j$	Number of trips for transportation of part j between cell i and i' in period h
$f_{ii'h}^j = \begin{cases} \left\lceil \frac{D_{jh}}{B_j} \right\rceil & \text{if } R_{i'j} - R_{ij} = 1 \\ 0 & \text{if } R_{i'j} - R_{ij} \neq 1 \end{cases}$	(1)
T_i	Capacity of machine i in each period
t_{ij}	Processing time of part j by machine i
SP	Set of (i, j) for parts assigned to machines
A_i	Fixed cost of machine i
B_i	Cost of increase/decrease of machine i which is considered to be fixed
α_j	The coefficient for each extra exception machine

Variables

$$\begin{aligned}
 X_{ikh} &= \begin{cases} 1 & \text{if machine } i \text{ is assigned to cell } k \text{ in period } h \\ 0 & \text{otherwise} \end{cases} \\
 Y_{jkh} &= \begin{cases} 1 & \text{if part } j \text{ is assigned to cell } k \text{ in period } h \\ 0 & \text{otherwise} \end{cases} \\
 Z_{iph} &= \begin{cases} 1 & \text{if machine } i \text{ in period } h \text{ is assigned to position } p \\ 0 & \text{otherwise} \end{cases} \\
 U_{ijkh} &= \begin{cases} 1 & \text{if } X_{ikh} = 1 \text{ and } Y_{jkh} = 0 \\ 0 & \text{otherwise} \end{cases}
 \end{aligned}$$

V_{ijkh}	$\begin{cases} 1 & \text{if } Y_{jkh} = 1 \text{ and } X_{ikh} = 0 \\ 0 & \text{otherwise} \end{cases}$
KA_{ih}	Number of machines type i which are added in each period h
KR_{ih}	Number of machines type i which are deleted in each period h
N_{ih}	Number of machines type i which are required for each period h

Therefore, the proposed model of this paper is stated as follows,

$$C_{ii'h}^j = \begin{cases} (M_{i'h} - M_{ih})C_{intra}^j & \text{if } X_{ikh}, X_{i'kh} > 0, M_{i'h} > M_{ih} \\ (M_{ih} - M_{i'h})C_{back}^j & \text{if } X_{ikh}, X_{i'kh} > 0, M_{i'h} < M_{ih} \\ |k - k'|C_{inter}^j & \text{if } X_{ikh}X_{i'kh} = 0, X_{ikh}X_{i'k'h} > 0 \end{cases} \quad (2)$$

$$\min \sum_{h=1}^H \sum_{j=1}^n \sum_{i=1}^m \sum_{i'=1}^m f_{ii'h}^j C_{ii'h}^j + \sum_{h=2}^H \sum_{i=1}^m B_i (KA_{ih} + KR_{ih}) + \sum_{h=1}^H \sum_{i=1}^m A_i \cdot N_{ih} + \sum_{h=1}^H \sum_k \sum_j \sum_{(i,j) \in sp} \alpha_j \cdot \frac{(U_{ijkh} + V_{ijkh})}{2} \quad (3)$$

subject to

$$\sum_{k=1}^C X_{ikh} = 1, \quad i = 1, 2, \dots, m, \quad \forall h \quad (4)$$

$$\sum_{k=1}^C Y_{jkh} = 1, \quad j = 1, 2, \dots, n, \quad \forall h \quad (5)$$

$$\sum_{i=1}^m X_{ikh} \leq NM, \quad k = 1, 2, \dots, C, \quad \forall h \quad (6)$$

$$\sum_{i=1}^m Z_{iph} = 1, \quad p = 1, 2, \dots, mp, \quad \forall h \quad (7)$$

$$\sum_{p=1}^{mp} Z_{iph} = 1, \quad i = 1, 2, \dots, m, \quad \forall h \quad (8)$$

$$\left(\sum_{j=1}^n D_{jh} \cdot t_{ij} \right) / T_i \leq N_{ih} \quad \forall i, h \quad (9)$$

$$N_{i(h-1)} + KA_{ih} - KR_{ih} = N_{ih} \quad \forall i, k, h \quad (10)$$

$$KA_{ih} \times KR_{ih} = 0 \quad \forall i, h \quad (11)$$

$$N_{ih}, KA_{ih}, KR_{ih} \geq 0 \quad \text{and} \quad \text{Integer}$$

$$X_{ikh}, Y_{jkh}, Z_{iph}, W_{kk'h}^i, U_{ijkh}, V_{ijkh} = 0 \text{ or } 1 \quad (12)$$

The proposed model given in Eqs. (1-12) presents a dynamic method to minimize different costs including the total cost of movements within and between cells and exceptional parts. The resulted problem is a nonlinear mixed integer programming problem. The objective function consists of four parts. The first part is associated with the cost of intra as well as inter cells transports. The second part of the objective function is associated with the cost of adding/removing a machine inside a cell. The third part of the objective function is associated with the fixed cost maintenance, depreciation, etc. Finally, the last part tries to minimize the number of exceptional parts.

Constraint (4) assures that each machine is assigned to only one cell. Eq. (5) specifies that each part is assigned only to one part group. Eq. (6) limits the number of machines assigned to each cell. Eq. (7) is to make sure that in each position only one machine is assigned. Eq. (8) is to make sure that in each period each machine is assigned to a particular position. Eq. (9) is to assure that in each period the load on each machine is not more than its capacity. Eq. (10) balances the number of machines, which are added/removed. According to Eq. (11), in each period, either a machine is added or removed. Finally, Eq. (12) determines the type of variables.

3. Solution procedure

The proposed study of this paper uses genetic algorithm to solve the results problem formulation.

3.1 Genetic Algorithm

Genetic algorithms are considered a method with an intelligent random search (Goldberg & Holland, 1988) where in each iteration, it attempts to evolve the current generation into a new population. This procedure is repeated using selection, crossover, and mutation operators. Some of the chromosomes in the current generation are chosen by selection mechanism and used for the next generation. Also, some individuals are chosen from current population, as parents, to setup offspring by implementing the crossover operator. Finally, mutation operator is implemented to change some genes in some chromosomes. In other words, the mutation operator makes sure about the diversity of searching in solution space. The chromosome in this paper is defined in Fig. 1 as follows,

X_1, X_2, \dots, X_m	Y_1, Y_2, \dots, Y_m	Z_1, Z_2, \dots, Z_n
The numbers for machines	The sequence of machines	Part allocation

Fig. 1. The arrangement of chromosome

According to Fig. 1, we understand that the longer the period of planning, the bigger the size of the chromosome. The length of the chromosome can be defined as follows,

$$(\text{Number of different types of machines} \times 2 + \text{Number of parts}) \times (\text{Number of planning programs})$$

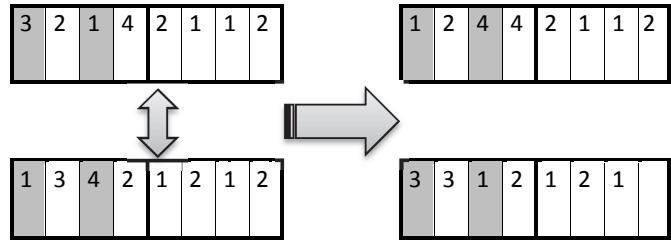
For instance, when there are 3 machines, 4 parts and 2 period of planning, the length of the chromosome will be 20 and we have

First period																			
Machine numbers			Allocation of machines to cells						Allocation of parts to cells										
2	1	3	2	1	1	2	1	1	2	2	3	1	2	1	1	2	2	1	1

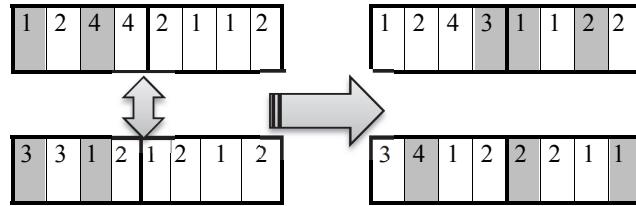
1	2	3	1	2	3	1	2	3	4	1	2	3	1	2	3	1	2	3	4
2	1	3	2	1	1	2	1	1	2	2	3	1	2	1	1	2	2	1	1

Fig. 2. The structure of Chromosome for 3 machines, 4 parts and 2 period of planning

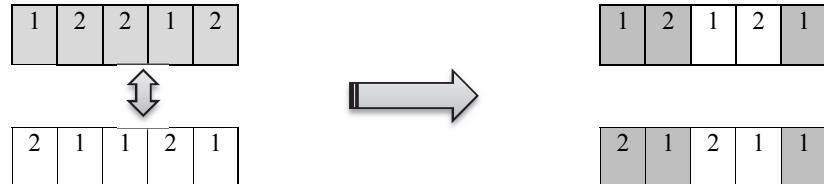
The proposed study normally generates 50-100 chromosome during the initiation of the proposed study. Next, we perform crossover on the operations as shown in Fig. 3 as follows,

**Fig. 3.** Discrete crossover

In Fig. 3, there are some violations of constraints. For instance, in the first chromosome, machine 4 is assigned to two positions 3 and 4 while machine 3 is assigned to no position. The chromosome is adjusted as shown in Fig. 4 as follows,

**Fig. 4.** The adjusted chromosome

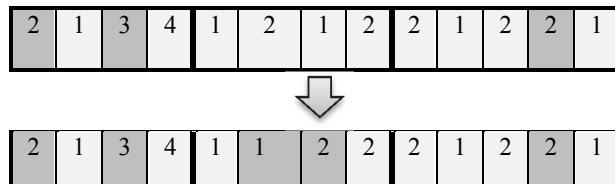
The third part of the chromosome, which is assigned for part family can be represented in Fig. 5 as follows,

**Fig. 5.** Discrete crossover to create an infant in the cell formation

After applying the crossover operator, chromosomes go for mutation operation. We create random number for each cell of the chromosome and in case this number is less than P_m the mutation is performed on that cell. In case the mutation operation is illegal according to the rules, we make some adjustments. For instance, consider the chromosome with 4 machines and 5 parts shown in Fig. 6 as follows,

**Fig. 6.** A chromosome with 4 machines and 5 parts

To adjust the chromosome after implementing the mutation we have Fig. 6 as follows,

**Fig. 7.** Discrete crossover to create an infant

The proposed study considers different conditions for terminating the algorithm such as reaching a predefined number of iterations, time, etc.

4. The results

In order to measure the performance of the proposed method, we have solved the proposed method for some instances using an optimization software and compared the results with the proposed GA. Table 1 demonstrates the input parameters used for the proposed study based on uniform distribution function.

Table 1

The input data used for the proposed study

Parameter	Value
B_j	$U(20, 30)$
D_{jh}	$U(200, 300)$
C_{intra}^j	1
C_{back}^j	3
C_{inter}^j	7
C_i	$U(3000, 4000)$
B_i	$U(4500, 5500)$
A_i	$U(13000, 17000)$
T_i	1000
t_{ij}	$U(1, 2)$
α_j	$U(80, 100)$

We have used these numbers to solve the problem under three different conditions.

4.1. Problem with 10 parts and 8 machines

For this problem, the population of each generation is set to 100 ($\mu = 100$), the probability for crossover and mutation are ($p_c = 0.9$) and ($p_m = 0.1$), respectively. In addition, termination criterion is to reach 100 iterations. In addition, we assume there are three types of machines in each cell, there are three cells, the average number of operations is equal to 3 and machine part is as follows,

	p_1	p_2	p_3	p_4	p_5	p_6	p_7	p_8	p_9	p_{10}
M_1	0	1	0	3	0	0	2	0	0	1
M_2	1	0	2	0	0	0	3	1	0	0
M_3	0	0	0	1	2	3	0	0	0	0
M_4	3	2	0	0	0	0	1	0	2	0
M_5	0	0	3	0	1	0	0	2	0	3
M_6	0	0	1	0	3	0	0	0	1	0
M_7	0	3	0	2	0	1	0	0	0	2
M_8	2	0	0	0	0	2	0	3	0	0

Table 2 presents details of the final solutions for two methods.

Table 2

The results of comparison of GA versus optimal solution using Lingo software

Obj. Func.	Rent	Add/remove cost	Displacement of machines	Displacement of parts	Exceptional parts	Time
Lingo	125959	125280	0	0	372	307 148.4
GA	125959	125280	0	0	372	307 32.47

According to the results of Table 2, both methods reach the optimal solution and with the same numbers.

4.2. Problem with 14 parts and 9 machines

For this problem, the population of each generation is set to 100 ($\mu = 100$), the probability for crossover and mutation are ($p_c = 0.9$) and ($p_m = 0.1$), respectively. Moreover, termination criterion is to reach 100 iterations. In addition, we assume there are three types of machines in each cell, there are three cells, the average number of operations is equal to 4 and machine part is as follows,

	p_1	p_2	p_3	p_4	p_5	p_6	p_7	p_8	p_9	p_{10}	p_{11}	p_{12}	p_{13}	p_{14}
M_1	0	1	0	3	0	0	2	0	0	1	0	0	2	0
M_2	1	0	2	0	0	0	3	1	0	0	0	2	0	0
M_3	0	0	0	1	2	4	0	0	0	0	3	0	3	1
M_4	3	2	0	0	0	0	1	0	2	0	0	1	0	0
M_5	0	0	3	0	1	0	0	2	0	3	0	0	0	3
M_6	0	0	1	0	3	0	0	0	1	0	2	0	0	4
M_7	0	3	0	2	0	1	0	4	0	4	0	3	0	0
M_8	2	0	0	0	4	2	0	3	0	0	1	0	0	2
M_9	0	0	4	0	0	3	0	0	3	2	0	0	1	0

Table 3 presents details of the final solutions for two methods.

Table 3

The results of comparison of GA versus optimal solution using Lingo software

	Obj. Func.	Rent	Add/remove cost	Displacement of machines	Displacement of parts	Exceptional parts	Time
Lingo	142242.5	140652	0	0	903	687.5	902.85
GA	142242.5	140652	0	0	903	687.5	50.51

One more time, both methods have reached the same results and they are different only in terms of CPU time. For this problem, the first cell contains machine number 4, 6 and 7 and parts, 1, 7 and 9. The second cell includes machines 1, 2 and 8 and parts {2, 5, 4, 6, 8, 10, 11, 13, 14}. Finally, cell number 3 includes {3, 5, 9} and parts {3, 12}.

4.3. Problem with 14 parts and 7 machines

For this problem, the population of each generation is set to 100 ($\mu = 100$), the probability for crossover and mutation are ($p_c = 0.9$) and ($p_m = 0.1$), respectively. Moreover, termination criterion is to reach 100 iterations. In addition, we assume there are three types of machines in each cell, there are three cells, the average number of operations is equal to 3, there are two periods and machine part is as follows,

	p_1	p_2	p_3	p_4	p_5	p_6	p_7	p_8	p_9	p_{10}	p_{11}	p_{12}	p_{13}	p_{14}
M_1	0	2	0	0	1	0	0	1	1	0	0	0	0	1
M_2	3	0	0	2	3	3	1	0	0	0	0	0	0	0
M_3	0	0	0	0	0	0	0	0	0	2	2	2	1	0
M_4	0	0	1	0	0	0	2	0	2	1	1	1	0	0
M_5	0	1	0	0	0	0	0	2	0	0	0	0	0	2
M_6	2	0	0	0	0	2	0	3	0	0	0	0	0	0
M_7	1	0	0	1	2	1	0	0	3	0	0	0	0	2

Table 4 shows details of the final solutions for two methods.

Table 4

The results of comparison of GA versus optimal solution using Lingo software

	Obj. Func.	Rent	Add/remove cost	Displacement of machines	Displacement of parts	Exceptional parts	Time
Lingo	147117.2	112994	20684	10750	1564.2	1125	2407.8
GA	147117.2	112994	20684	10750	1564.2	1125	131.02

One more time, both methods reached the optimal solutions under different CPU time. The following demonstrate the position of machines to different locations.

	Position 1	Position 2	Position 3	Position 4	Position 5	Position 6	Position 7
Period 1	3	4	7	2	1	5	6
Period 2	3	4	1	7	5	6	2

In addition, the positions of machines to different cells.

	M_1	M_2	M_3	M_4	M_5	M_6	M_7
Period 1	2	2	1	1	3	3	2
Period 2	2	3	1	1	3	3	2

Finally, the allocation of parts to cells are summarized as follows,

	p_1	p_2	p_3	p_4	p_5	p_6	p_7	p_8	p_9	p_{10}	p_{11}	p_{12}	p_{13}	p_{14}
Period 1	3	1	3	1	2	1	1	3	3	3	3	3	2	1
Period 2	3	1	3	3	2	3	2	2	1	2	3	1	3	2

4. Conclusion

In this paper, we have presented a dynamic method to minimize different costs including the total cost of movements within and between cells and exceptional parts. In this study, the Hierarchical Genetic Algorithm (HGA) has been used for solving the resulted model and the results are compared with genetic algorithm. The results have indicated that the proposed model was capable of detecting optimal solutions in reasonable amount of time for small and middle scale problems. This survey can be extended to a new research areas by considering uncertainty on different parameters such as processing times, cost of transportation, etc. and we leave it as a future research for interested researchers.

Acknowledgement

The authors would like to thank the anonymous referees for constructive comments on earlier version of this paper.

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